

Supplementary Material of P-Tucker

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Abstract

In this supplementary material, we suggest full proof of the row-wise update rule for factor matrices and theoretical complexities of P-TUCKER-APPROX, which were introduced in the main paper.

I. FULL PROOF OF THE ROW-WISE UPDATE RULE

Definition 1 (Sparse Tucker Factorization): Given a tensor \mathcal{X} ($\in \mathbb{R}^{I_1 \times \dots \times I_N}$) with observable entries Ω , the goal of sparse Tucker factorization of \mathcal{X} is to find factor matrices $\mathbf{A}^{(n)}$ ($\in \mathbb{R}^{I_n \times J_n}$) and a core tensor \mathcal{G} ($\in \mathbb{R}^{J_1 \times \dots \times J_N}$), which minimize Equation (1).

$$L(\mathcal{G}, \mathbf{A}^{(1)}, \dots, \mathbf{A}^{(N)}) = \sum_{\forall (i_1, \dots, i_N) \in \Omega} \left(\mathbf{x}_{(i_1, \dots, i_N)} - \sum_{\forall (j_1, \dots, j_N) \in \mathcal{G}} \mathcal{G}_{(j_1, \dots, j_N)} \prod_{n=1}^N a_{i_n j_n}^{(n)} \right)^2 + \lambda \sum_{n=1}^N \|\mathbf{A}^{(n)}\|^2 \quad (1)$$

Theorem 1 (Row-wise Update rule for Factor Matrices):

$$\arg \min_{[a_{i_n 1}^{(n)}, \dots, a_{i_n J_n}^{(n)}]} L(\mathcal{G}, \mathbf{A}^{(1)}, \dots, \mathbf{A}^{(N)}) = \mathbf{c}_{i_n}^{(n)} \times [\mathbf{B}_{i_n}^{(n)} + \lambda \mathbf{I}_{J_n}]^{-1} \quad (2)$$

$$\text{where the } (j_1, j_2)\text{th entry of } \mathbf{B}_{i_n}^{(n)} (\in \mathbb{R}^{J_n \times J_n}) : \sum_{\forall (i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \delta_{(i_1, \dots, i_N)}^{(n)}(j_1) \delta_{(i_1, \dots, i_N)}^{(n)}(j_2), \quad (3)$$

$$\text{the } j\text{th entry of } \mathbf{c}_{i_n}^{(n)} (\in \mathbb{R}^{J_n}) : \sum_{\forall (i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \mathbf{x}_{(i_1, \dots, i_N)} \delta_{(i_1, \dots, i_N)}^{(n)}(j), \quad (4)$$

$$\text{the } j\text{th entry of } \delta_{(i_1, \dots, i_N)}^{(n)} (\in \mathbb{R}^{J_n}) : \sum_{\forall (j_1 \dots j_{n-1}, j, j_{n+1} \dots j_N) \in \mathcal{G}} \mathcal{G}_{(j_1 \dots j_{n-1}, j, j_{n+1} \dots j_N)} \prod_{k \neq n} a_{i_k j_k}^{(k)}. \quad (5)$$

Proof:

$$\frac{\partial L}{\partial a_{i_n j_n}^{(n)}} = 0, \forall j_n, 1 \leq j_n \leq J_n$$

$$\Leftrightarrow \sum_{\forall (i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \left(\left(\mathbf{x}_{(i_1, \dots, i_N)} - \sum_{\forall (j_1, \dots, j_N) \in \mathcal{G}} \mathcal{G}_{(j_1, \dots, j_N)} \prod_{n=1}^N a_{i_n j_n}^{(n)} \right) \times \left(-\delta_{(i_1, \dots, i_N)}^{(n)}(j_n) \right) \right) + \lambda a_{i_n j_n}^{(n)} = 0$$

$$\Leftrightarrow \sum_{\forall (i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \left(\left(\sum_{t=1}^{J_n} \delta_{(i_1, \dots, i_N)}^{(n)}(t) a_{i_n t}^{(n)} \right) \times \left(\delta_{(i_1, \dots, i_N)}^{(n)}(j_n) \right) \right) + \lambda a_{i_n j_n}^{(n)} = \sum_{\forall (i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \left(\mathbf{x}_{(i_1, \dots, i_N)} \delta_{(i_1, \dots, i_N)}^{(n)}(j_n) \right) \quad (6)$$

$\sum_{\forall(i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \left(\left(\sum_{t=1}^{J_n} \delta_{(i_1, \dots, i_N)}^{(n)}(t) a_{i_n t}^{(n)} \right) \times \left(\delta_{(i_1, \dots, i_N)}^{(n)}(j_n) \right) \right)$ is expressed as an inner product of the following vectors.

Row vector ($1 \times J_n$) : $[a_{i_n 1}^{(n)}, \dots, a_{i_n J_n}^{(n)}]$

Column vector ($J_n \times 1$) : $\begin{bmatrix} \sum_{\forall(i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \delta_{(i_1, \dots, i_N)}^{(n)}(1) \delta_{(i_1, \dots, i_N)}^{(n)}(j_n) \\ \vdots \\ \sum_{\forall(i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \delta_{(i_1, \dots, i_N)}^{(n)}(J_n) \delta_{(i_1, \dots, i_N)}^{(n)}(j_n) \end{bmatrix}$

If we vary j_n from 1 to J_n , the row vector is fixed as $a_{i_n \cdot}^{(n)}$ and the column vector differs. Thus, we can integrate each column vector into a matrix $\mathbf{B}_{i_n}^{(n)} (\in \mathbb{R}^{J_n \times J_n})$

where the (j_1, j_2) th entry of $\mathbf{B}_{i_n}^{(n)} (\in \mathbb{R}^{J_n \times J_n})$: $\sum_{\forall(i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \delta_{(i_1, \dots, i_N)}^{(n)}(j_1) \delta_{(i_1, \dots, i_N)}^{(n)}(j_2)$.

λ term is simply transformed into $\lambda \mathbf{I}_{J_n}$, where \mathbf{I}_{J_n} is an identity matrix ($\in \mathbb{R}^{J_n \times J_n}$). In the same way, the right part of Equation (6) is integrated as $\mathbf{c}_{i_n \cdot}^{(n)} (\in \mathbb{R}^{J_n})$

where the j th entry of $\mathbf{c}_{i_n \cdot}^{(n)} (\in \mathbb{R}^{J_n})$: $\sum_{\forall(i_1, \dots, i_N) \in \Omega_{i_n}^{(n)}} \mathbf{x}_{(i_1, \dots, i_N)} \delta_{(i_1, \dots, i_N)}^{(n)}(j)$.

Therefore, Equation (6) is equivalent to

$$[a_{i_n 1}^{(n)}, \dots, a_{i_n J_n}^{(n)}] \times [\mathbf{B}_{i_n}^{(n)} + \lambda \mathbf{I}_{J_n}] = \mathbf{c}_{i_n \cdot}^{(n)}$$

Since $\mathbf{B}_{i_n}^{(n)}$ is represented as the sum of rank-1 matrices and $\lambda > 0$, matrix $[\mathbf{B}_{i_n}^{(n)} + \lambda \mathbf{I}_{J_n}]$ is positive-definite and invertible. Hence,

$$\Leftrightarrow [a_{i_n 1}^{(n)}, \dots, a_{i_n J_n}^{(n)}] = \mathbf{c}_{i_n \cdot}^{(n)} \times [\mathbf{B}_{i_n}^{(n)} + \lambda \mathbf{I}_{J_n}]^{-1}$$

■

Algorithm 1: P-TUCKER for Sparse Tensors

Input : Tensor $\mathcal{X} \in \mathbb{R}^{I_1 \times I_2 \times \dots \times I_N}$,
core tensor dimensionality J_1, \dots, J_N , and
truncation rate p (P-TUCKER-APPROX only).
Output: Updated factor matrices $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times J_n}$ ($n = 1, \dots, N$),
and updated core tensor $\mathcal{G} \in \mathbb{R}^{J_1 \times J_2 \times \dots \times J_N}$.

- 1 initialize factor matrices $\mathbf{A}^{(n)}$ ($n = 1, \dots, N$) and core tensor \mathcal{G}
- 2 **repeat**
- 3 update factor matrices $\mathbf{A}^{(n)}$ ($n = 1, \dots, N$)
- 4 calculate reconstruction error
- 5 **if** P-TUCKER-APPROX **then** ▷ \mathcal{G} Truncation
- 6 remove “noisy” entries of \mathcal{G} by Algorithm 2
- 7 **until** the maximum iteration or $\|\mathcal{X} - \mathcal{X}'\|$ converges;
- 8 **for** $n = 1 \dots N$ **do**
- 9 $\mathbf{A}^{(n)} \rightarrow \mathbf{Q}^{(n)} \mathbf{R}^{(n)}$ ▷ QR decomposition
- 10 $\mathbf{A}^{(n)} \leftarrow \mathbf{Q}^{(n)}$ ▷ Orthogonalize $\mathbf{A}^{(n)}$
- 11 $\mathcal{G} \leftarrow \mathcal{G} \times_n \mathbf{R}^{(n)}$ ▷ Update core tensor \mathcal{G}

Algorithm 2: P-TUCKER-APPROX

Input : Tensor $\mathcal{X} \in \mathbb{R}^{I_1 \times I_2 \times \dots \times I_N}$,
factor matrices $\mathbf{A}^{(n)} \in \mathbb{R}^{I_n \times J_n}$ ($n = 1, \dots, N$),
core tensor $\mathcal{G} \in \mathbb{R}^{J_1 \times J_2 \times \dots \times J_N}$, and
truncation rate p ($0 < p < 1$).
Output: Truncated core tensor $\mathcal{G}' \in \mathbb{R}^{J_1 \times J_2 \times \dots \times J_N}$.

- 1 **for** $\beta = \forall(j_1, \dots, j_N) \in \mathcal{G}$ **do**
- 2 compute a partial reconstruction error $\mathcal{R}(\beta)$
- 3 sort $\mathcal{R}(\beta)$ in a descending order with their indices
- 4 remove $p|\mathcal{G}|$ entries in \mathcal{G} , whose $\mathcal{R}(\beta)$ value are ranked within top- $p|\mathcal{G}|$ among all $\mathcal{R}(\beta)$ values.

II. THEORETICAL COMPLEXITIES OF P-TUCKER-APPROX

Theorem 2 (Time complexity of P-TUCKER-APPROX): The time complexity of P-TUCKER-APPROX is $O(NIJ^3 + N^2|\Omega||\mathcal{G}|)$.

Proof: The only difference between P-TUCKER and P-TUCKER-APPROX is that P-TUCKER-APPROX exploits $|\mathcal{G}|$ entries rather than using full J^N entries of \mathcal{G} . Thus, the time complexity of P-TUCKER-APPROX for updating factor matrices and computing the reconstruction error is reduced to $O(NIJ^3 + N^2|\Omega||\mathcal{G}|)$. Moreover, the cost of Algorithm 2 is $O(N|\Omega||\mathcal{G}|)$, which is much less than that of other parts. Hence, the time complexity of P-TUCKER-APPROX is $O(NIJ^3 + N^2|\Omega||\mathcal{G}|)$. ■

Theorem 3 (Memory complexity of P-TUCKER-APPROX): The memory complexity of P-TUCKER-APPROX is $O(J^N)$.

Proof: Compared to P-TUCKER, P-TUCKER-APPROX requires additional intermediate data to store $\mathcal{R}(\beta)$. The memory complexity of $\mathcal{R}(\beta)$ is at most $O(J^N)$, and the memory requirements for $\mathcal{R}(\beta)$ is much larger than other intermediate data. Therefore, the memory complexity of P-TUCKER-APPROX is $O(J^N)$. ■